

A COMPREHENSIVE ANALYSIS ON VARIOUS ADVANCED WEIGHT BASED CLUSTERING METHODOLOGIES IN MANET

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ABSTRACT

Mobile ad-hoc network is an explicit sort of remote network that can be immediately conveyed without prior framework. The qualities of MANET are a dynamic topology, dispersed activity, multihop and so forth. Clustering is an essential research territory; it offers a few advantages like it enhances strength and diminishes the overhead of the network that expands the effectiveness of the network. There are many clustering plans have been proposed for MANETs. An orderly grouping of these clustering plans empowers one to all the more likely comprehend and make an enhancement. Each clustering algorithm utilizes different parameter for choosing group head in the bunch. The group head is fills in as a leader in bunch and keep up the entire network data. A similar investigation of different methods additionally has been canvassed in this paper. This paper introduces a study of various clustering plans.

Keywords: Mobile Ad-Hoc Network, Clustering, Cluster Head, Cluster Gateway, Cluster Member.

I. INTRODUCTION

Alongside the advancement of web and correspondence aptitude, MANETs have pulled in wide research exertion as of late, before, numerous methodologies were proposed to complete the issue coming in way steering, node clustering, and numerous others in MANETs. A mobile ad-hoc network is a gathering of the network which all things considered structures a network as we require. In MANET all nodes are mobile i.e. they every now and again change their position the all nodes are associated with one another through a remote medium. To send an information bundle in a network they utilize a directing protocol. Directing protocol gives a course from source to goal with the goal that information can be sent effortlessly.

In MANET there are three steering protocol, Proactive, Reactive, Hybrid. In Proactive steering protocol or table-driven directing protocol, every node keeps up the steering table in which information of each node and their way is kept up, at whatever point any topology of the node is transformed it is promptly modernized in the table and quickened to all node. With the goal that each node must have a refreshed table and each node is capable if any change is done it must be refreshed in a directing table. A portion of the proactive steering protocols are DSDV, WRP. In Reactive directing protocol or on-request steering protocol, there is no steering table when a source node needs to interconnect with another node in a network the course disclosure process is activated. The course support is utilized to keep up the course and discover the blunder between the courses. A portion of the responsive protocols are AODV, DSR. In half breed directing protocol it incorporates the property of both responsive and additionally the proactive steering protocol. ZRP is one of the mixtures directing protocol.

In directing in a network incorporates a substantial number of nodes which is a noteworthy issue. To conquer this overhead we use clustering.

II. Clustering Types And Its Advantages:-

In a cluster, nodes are gathered by their area and load in a mobile ad-hoc network. As we make these little gatherings from an outsized number of nodes so they are progressively steady and responsive moreover. The nodes in a cluster build up their association utilizing broadcast message. At the point when nodes in cluster move from one cluster to other then it expected to refresh the information, as the node lives in a cluster there is no compelling reason to change the whole network.

The cluster head, gateway, and nodes in a cluster are the real parts which drive a critical part in clustering; Cluster head and gateway are spines of the design of clustering. There is a node in a cluster which goes about as a cluster head and different nodes are treated as a cluster member. There is a system to choose a cluster head. Cluster gateway is a typical node between two clusters which neither a cluster head nor a cluster member. CH gives a connection between two clusters which exchanges information from one cluster to other.

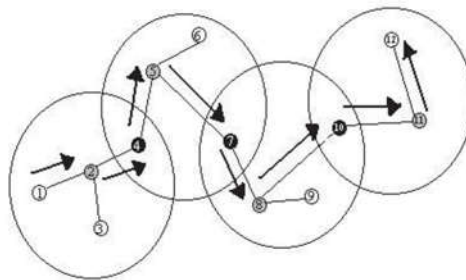


Fig 1:- (Cluster Structure in MANET) [4].

In a cluster, there are two sorts of correspondence, Intra-cluster correspondence, and Inter-cluster correspondence.

In Intra-cluster correspondence, the correspondence is 1-bounce correspondence with each cluster member. Consequently, the cluster head will straightforwardly interconnect with member nodes. Yet, member nodes can't specifically speak with different nodes of the cluster.

In Inter-cluster correspondence, the correspondence is a multipoint hand-off which chooses the cluster head through which information bundle is exchanged. Because of this it diminishes the quantity of nodes engaged with sending the information parcel and decreases the use of battery and expands the proficiency.

2.1 Types Of Clustering [4]:-

- DS-Based clustering: Finding an overwhelming set to diminish the quantity of nodes taking part course seek. Ex Connected DS.
- Low-Maintenance clustering: Providing a cluster framework for upper layer application. Ex Least cluster change.

- Mobile-Aware clustering: Utilizing mobile nodes portability nature for cluster arrangement and its consideration. Ex MOBIC.
- Energy-Efficient clustering: It maintains a strategic distance from the superfluous usage of vitality and gives better utilization of vitality. Ex-Energy based DS.

III. CLUSTERING ALGORITHMS FOR MANET

In contrast to traditional remote networks, the ad-hoc network is a framework less Network comprising of mobile self-sufficient moving nodes. Every node in the network goes about as a switch or a bundle forwarder and they interconnect with one another at a similar friend level to empower the network work. While steering through nodes a portion of the directing issues will emerge. To conquer such issues in MANET directing, clustering system has been utilized.

Clustering approach is a critical research point for MANETs and generally utilized in productive network the board, progressive steering protocol configuration, network demonstrating,

Nature of Service. Numerous scientists' ongoing spotlight has been on clustering the board that is one of the major issues in MANETs. The principle target of clustering in mobile ad hoc network conditions is proficient cluster head choice and enhancement of QoS.

Using the clustering algorithm to develop progressive topology might be a decent strategy to take care of these issues. An adaptive mobile cluster algorithm can support the portability impeccably and keeps up the steadiness and power of network engineering. To help the multi-jump and mobile qualities of remote ad hoc network, the quick organization of network and dynamic reproduction after topology changes are viably executed by clustering the board.

Clustering the board has five particular advantages over different protocols. Right off the bat, it utilizes different channels viably and enhances framework limit significantly Secondly, it lessens the trade overhead of control messages and fortifies node the board [16] thirdly; it is anything but difficult to execute the nearby synchronization of the network. It gives nature of administration (QoS) directing administration productively. At last, it can bolster the remote networks with an extensive number of nodes.

Cluster-based directing is one of the steering plans for MANET in which different clusters of mobile nodes are shaped with each gathering called a cluster. These nodes sort out themselves around another node through cluster head. The cluster head goes about as a nearby organizer for its cluster to perform between cluster steering, information sending, etc Clustering of nodes spares vitality and correspondence data transfer capacity in ad-hoc networks.

A considerable measure of research has been done on clustering algorithm for MANETs. The accompanying involves overview of different clustering algorithms and their exhibitions are investigated utilizing different execution measurements and its similar examination.

IV. Mobility Based Clustering Algorithm

. In this methodology, portability of the node is mulled over for cluster arrangement that accomplishes most extreme steadiness of collection the mobile nodes of comparative moving examples. A portion of the models are MOBIC, MobDhop, and so forth.

4.1 Mobility Based Clustering Algorithm (MOBIC)

A Lowest Relative Mobility Clustering Algorithm-MOBIC is like the Lowest-ID algorithm; however it utilizes the versatility metric as a premise of cluster development and cluster head determination. Instead of node ID, MOBIC utilizes portability metric, Aggregate Local Mobility (ALM) to choose a cluster head. Right off the bat, the mobile nodes dispersed over. The proportion between the got power dimensions of progressive transmissions between couples of nodes is utilized to process the relative portability between neighboring nodes, which decides the ALM of every node. Furthermore, every node broadcasts the portability metric of every node to its neighbors. Every node contrasts the versatility and they got total portability esteem and choose the node as a cluster head that has the most minimal versatility estimation of among every one of its neighbors.

On the off chance that the cluster has two cluster heads as a neighbor, one of the cluster heads changes the status as cluster gateway. Resort to an examination of node IDs and pursue a Lowest algorithm, if two cluster undecided state nodes or cluster heads that are in conflict to hold the cluster head state have a similar estimation of Mobility. At last, if two cluster heads came into a similar transmission, extend re-clustering is activated.

4.2 Mobility-based d-hop Clustering Algorithm (MobDHop)

In MobDhop, the whole network is apportioned into d-bounce clusters dependent on portability metric and the distance across of a cluster as for node versatility. The cluster head decision is reliant on three portability measurements, specifically (i) Variation of assessed separate between nodes after some time

that evaluates the general versatility of two nodes. (ii) Local Variability that computes the mean of the considerable number of neighbors for node (iii) Group changeability, which ascertains the mean of its 1-bounce neighbors.

A node that has the most reduced inconstancy esteem expects the job of cluster head. Pronounce the mobile node as cluster head, if its neighbor list is unfilled. Look at the nearby inconstancy of a mobile node with all its non-clustered neighbors; if its neighbor list isn't void choose the mobile node as cluster head that has most minimal neighborhood fluctuation esteem in respect to its neighbors. Non-clustered node solicitations to join its neighboring cluster. A node may move toward becoming non-clustered when it is recently actuated or it loses its cluster head because of node versatility.

4.3 Flooding Based Clustering Algorithm

The clustering depends on the recognized component of conveying information or control messages to all nodes inside the network for example, Max-Min D cluster algorithm.

Max-Min d-cluster Formation Algorithm [5]

The primary target of this algorithm is to sum up the cluster definition to a gathering of nodes that are up to d-jumps from a cluster head. Because of a substantial number of nodes contribution, it is attractive to give the nodes a chance to work no concurrently. The clock synchronization overhead is maintained a strategic distance from [38], giving additional handling investment funds. Moreover, the quantity of messages sent from every node is restricted to several of d the most extreme number of jumps from the closest cluster head, as opposed to n the quantity of nodes in the network. This ensures an all around controlled message unpredictability for the algorithm. Additionally, as d is an info incentive to the heuristic, there is a command over the quantity of cluster heads chose or the thickness of

cluster heads in the network. The measure of assets required at every node is insignificant, comprising of four straightforward principles and two information structures that keep up node data over 2d rounds of correspondence.

Nodes are contender to be clustered dependent on their node id instead of their level of network. As the network topology changes somewhat the node's level of availability is considerably more liable to change than the node's id with respect to its neighboring nodes. On the off chance that a node an is the biggest in the d-neighborhood of another node B then node An, A will be chosen a cluster head, despite the fact that node A may not be the biggest in its d-neighborhood. This gives a smooth trade of cluster heads as opposed to an inconsistent trade. This technique limits the measure of information that must be passed from an active cluster head to another cluster head when there exchange.

Cluster Based Algorithm

It utilizes half breed steering protocol and the clustering structure to diminish directing control overhead and enhance the adaptability of the network. In Cluster Based Routing Protocol (CBRP), a cluster head change is insignificant. For keeping up the cluster following two cluster head, change rules are forced in CBRP as portrayed in [41]. (I) a non-cluster head never difficulties the status of a current cluster head. (ii). At the point when two cluster heads move alongside one another over some stretch of time (Conflict Time like a flash) one of them will lose its job of the cluster head.

Thus, at whatever point a cluster head hears HELLO messages from another cluster head demonstrating a bidirectional connection builds up; it sets Cluster Time to terminate in Conflict Time. At the point when Cluster Time lapses, it will check on the off chance that it is still in struggle with the other cluster head, by checking if the other cluster head is still in its neighbor table. Assuming this is the case, it contrasts its very own ID and that of the other cluster head's. The one with a littler ID will keep on going about as a cluster head. The one with a greater ID surrenders its job as a cluster head and changes from Cluster Head to Cluster Member in its resulting HELLO messages. This may trigger a redesign of different clusters.

These principles promise a type of cluster soundness by deferring the cluster head change by Conflict Time after happening to two cluster heads in one another's range.

This keeps away from pointless cluster head change if their progressing time is not exactly or equivalent to Conflict Time yet on the off chance that progressing time is more the cluster head change is constrained.

V. Weight Based Clustering Algorithm

In the load based clustering approach the heaviness of the node is considered for choosing a cluster head for shaping stable clusters. In view of the writing overview the load based clustering algorithm gives the answer for QoS of the network. A portion of the Weight Based Clustering Algorithms (WBCA) is examined underneath.

6.1 Entropy-Based Weighted Clustering Algorithm (EBWCA)

The real inadequacy of WCA is visit reaffiliation that expands the network overhead. EBWCA conquers the downside of WCA and structures a steady network. This algorithm utilizes an entropy-based model for assessing the course steadiness in the ad-hoc network and choosing cluster head. EBWCA enhances the execution, especially on the quantity of clusters and the reaffiliation recurrence. Entropy presents vulnerability and is a proportion of the confusion in a framework. Accordingly, it is a superior marker of the dependability and versatility of the ad hoc network.

6.2 Vote Based Clustering Algorithm (VBC)

It depends on two parameters. One is leftover battery life and the other is mobile nodes. Every mobile node has a one of a kind character number. The data will be transmitted through the regular channel by sending hi message. In this algorithm, they presented another idea called vote. It comprises of the data about the area and power utilization of the nodes. Vote esteem is mobile nodes esteem, i.e. weighted aggregate of substantial neighbors and lingering power.

Every node sends a welcome message haphazardly amid a welcome cycle. Amid the main cycle, the node does not know whether it has a clustered or not. Be that as it may, in the second cycle, every node sends another welcome message in which the casting a ballot thing is set to its own vote esteem. Presently every node knows the sender with the most astounding vote and not has a place with a current cluster is its cluster head. In the following sending hi message, it sends the cluster head ID esteem moreover. On the off chance that more than one node gets a similar number of hi message, the node that has most reduced ID will go about as a cluster head. By proceeding with this methodology, each node knows its cluster head after two-hi message period.

6.3 Weight Based Adaptive Clustering Algorithm (WBACA)

In this algorithm, the clustering algorithm depends on the accessibility of position data by means of the Global Positioning System. It considers the accompanying parameters for the choice of cluster head: transmission control, transmission rate, portability of the node, degree and battery control. In this algorithm, the node with the least weight is chosen as a cluster head. The rest of the nodes are the members of the cluster. The principle advantage of this algorithm is that it doesn't permit cluster heads to be one-jump neighbors of one another. Covering clusters are associated through Gateways. The node that interfaces two cluster head is called as a gateway node. Every other node are a member in one-bounce from their cluster head.

6.4 Maximal Weighted Independent Set in Wireless Networks (MWIWN)

The primary target of this algorithm is to locate a maximal weighted autonomous arrangement of nodes in a network. Every node is allocated a load dependent on its appropriateness of being a cluster head, which is registered considering the speed of the node. The node that has maximal weight is chosen as a cluster head. Slower nodes will have a greater load than the quicker ones. On account of MANET, parceling the node into gatherings (clusters) is essential for limiting the measure of information to be traded so as to keep up steering to control mobile condition. Be that as it may, it has been demonstrated that it gives more arrangement, for example, limiting the measure of capacity for correspondence data, conveying assets all through the network, and so forth.

MWIS has numerous advantages: it requires just the information of the nearby topology at every node, it is quick and simple to execute, and its time multifaceted nature is ended up being bound by a topology-subordinate parameter of the network, the strength number. The shortcoming of this algorithm is that it expects that the message sent by a node is gotten effectively inside a limited time by the entirety of its neighbors. This implies, a node needs to sit tight for every one of the reactions from its neighbors to settle on its own choice to be a cluster head or a standard node.

6.5 A Load-Balancing Weighted Clustering Algorithm in MANET

In this algorithm, the cluster head determination is adaptive, in light of the portability of the node or change in the relative separation dependent on nodes and the cluster head. It expects a predefined number of nodes that a cluster could cover. At the point when a cluster measure surpasses its predefined limit, new race methodology begins to adjust the quantity of nodes in the cluster. For better correspondence, the separation between the cluster head and cluster member is inside the transmission run. It is on the grounds that the relative separation between nodes influences the utilization of the battery control. Thus, the cluster head have additional duty to send bundles to different nodes. Versatility is a standout amongst the most critical difficulties in MANET design since it would change the network topology. A proficient cluster head ought not to move rapidly, on the grounds that when the cluster head changes quick, the nodes might be moved out of a cluster and are joined to another current cluster and in this manner bring about decreasing the soundness of the network. The primary point of this algorithm is to enhance the strength and enhance the load adjusting.

6.6 Efficient Management Algorithm for Clustering (EMAC)

EMAC [24] is another dispersed clustering algorithm where the node weight is determined utilizing components, for example, the node degree, remaining battery control, transmission power, and node versatility for the cluster heads decision. The principle point of this algorithm is to limit the quantity of clusters created. The algorithm utilizes the aggregate of separation utilized in WCA for transmission go that can cover the biggest range. In addition that the creator contended that residual battery control is a superior measure than the total time amid which the node goes about as a cluster head that is utilized in Weighted Clustering Algorithm since it permits expanding the lifetime of nodes by surrendering the job as a cluster head if there should be an occurrence of inadequate battery control. The algorithm likewise confines the quantity of nodes inside a cluster. By confining the quantity of nodes cooked by a cluster head helps in appropriate MAC working. The algorithm depends on the limit of the cluster and it utilizes the connection lifetime instead of the node versatility for the upkeep strategy.

The reason is the node versatility metric does not influence the decision of a CH as much as the connection strength metric does. EMAC executes diverse instruments for the entry of new nodes, cluster head nodes, member nodes, converging of clusters re-appointment of cluster heads.

6.7 Weighted Clustering Algorithm (WCA)

This methodology can progressively adapt itself with the regularly changing topology of ad hoc networks. Every node is relegated with a proper weight dependent on its reasonableness of being a cluster head. It chooses a cluster head as indicated by the quantity of nodes it can deal with, the portability of a node, transmission power, and battery control. To stay away from correspondences overhead, this algorithm isn't intermittent and the cluster head decision methodology is possibly summoned dependent on node portability and when the present overwhelming set is unable to cover every one of the nodes. To guarantee that cluster heads ought not to be overloaded, a pre-characterized edge is utilized which demonstrates the quantity of nodes that each cluster head can in a perfect world help. The node with the base weight is chosen as a cluster head. Once the cluster head decision is finished, every one of the nodes turn out to be either a cluster head or a member of a cluster head. The separation between members of a cluster head must be less or equivalent to the transmission go between them. No two cluster heads can be quick neighbors. The primary disadvantage of this algorithm is it devours more vitality in light of the fact that the cluster head has more obligations so it expends more vitality a thus visit re-alliance happens.

6.8 On-Demand Weighted Clustering Algorithm (ODWCA)

The primary point of ODWCA is to keep up the security of the network. Here the load is determined with the consolidated aftereffect of a few parameters, for example, perfect node degree, degree distinction, transmission control, cluster head serving time and versatility. Choosing the littlest weighted node as cluster head and limiting its neighbors to take an interest in further race technique is pursued. The fundamental advantage of this clustering plan is the adaptability of adjusting the weighting factor for every framework parameter to make it an appropriate situation.

CONCLUSION

In this paper, the overview of different weight based clustering plans has been finished. The strategies adopted in finding the cluster head are distinctive in these algorithms. This paper shows a review of clustering algorithms in which various measurements have been utilized to discover cluster head dependent on the load estimations of the node with other essential networks or parameters are mulled over. I trust that the overview exhibited in this paper will be useful and give scientists a stage for picking the correct weight based clustering algorithm for their work later on.

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